# Main goals of the assignment

- Learn about the compiler directives for multi-core system for computational-intensive tasks;
- Understand the differences between global and thread-private variables;
- Learn about the problem dimensions that are requiring parallel computations;
- Learn how to do a performance study for a parallel algorithm implementation.

## The problem to solve

#### General overview

A parallel version of matrix to matrix multiplication that is efficient is requested.

### Background

Let A and B be to  $n \times n$  matrices of real numbers. The elements of the product of the two matrices, noted by C in what follows is computed by

$$c_{ij} = \sum_{k=0}^{n-1} a_{ik} b_{kj}, \text{ for } 0 \le i, j < n$$
(1)

The time complexity of this multiplication is  $O(n^3)$ . When n is large, then the computational time is very high.

#### How to parallelize

See the slide for Lecture 4 for some examples. We will take the case in which we have p processing elements that are available (e.g. the number of our desktop cores).

In this exercise, we will split the matrix A in horizontal slices and we will assign one slice or more to one core. The core  $m \ (0 \le m < p)$  will therefore deal with the rows  $m \cdot \frac{n}{p} \le i < (m+1) \cdot \frac{n}{p}$  of A, but also of C. Note that is no overlaps of readings/writing from/to the global memories as was the case in Lab 1). (Be carefull to the use of loop variables!)

### To do

- 1. Write the sequential code to multiply the two matrices. The dimension of the matrices should be a given as parameter in the command line (hint: argv[1]). Dynamic allocation should be used (hint: see how we allocate the space for the vectors in lab 1). The two matrices' elements should be set to values that allows the simple check of the results (hint:  $a_{ij} = b_{ij} = 1$ ).
- 2. Write the parallel code to multiply the two matrices (hints: use the pragma for the external loop; be carefull how the loop variables are shared/or not). The number of cores should be a parameter in the command line (hint: argv[2]). Check the correctness of the result in simple cases.
- 3. Introduce time records (hint: omp\_get\_wtime) before and after the part that is parallelized.
- 4. Record the times  $T_p^{(n)}$  in table like the following (with the maximum cores that you have, e.g. 8 or 16) and compute the speedup
- 5. Compute the speedups using

$$S_p^{(n)} = rac{T_1^{(n)}}{T_n^{(n)}}$$

and record them in a similar table with the above one. Due the same for the efficiency  $E_p^{(n)} = S_p^{(n)}/p$ .

Table 1: Put inside the boxes the recorded times

$n \backslash p$	1	2	4	8
1600				
2000				
2400				

- 6. Display in a graphic the values of S as dependence on p (respectively E in another graphic) and with different polygonal lines (and colors) the values for different n
- 7. Draw conclusions related to:
  - increase/decrease of S with p;
  - $\bullet\,$  dependence of the problem dimension n